

Operating Systems for Sensor Networks

Michael Schulze

mschulze@ovgu.de

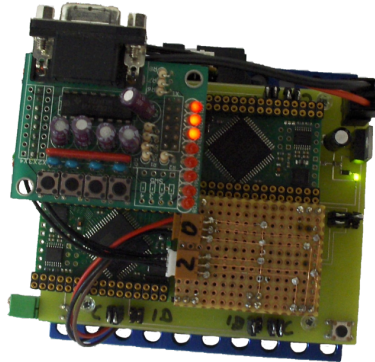
Department of Distributed Systems (IVS)
Embedded Systems and Operating Systems (EOS)
Otto-von-Guericke-Universität Magdeburg

Magdeburg, 09.07.2008



Outline

- Motivation
- TinyOS
- Contiki OS
- MANTIS OS
- Comparison
- Conclusion



Sensor Network Applications

- Environmental monitoring
 - Tracking of animals
 - Inventory control (RFID)
 - Home/office monitoring and automation
 - Structural monitoring
- Defense industry



Hardware

- Heterogeneity
- Miniaturization
- Integration
- Energy management

Typical Hardware of a sensor node e.g. Chipcon CC2420DB

- AVR ATmega128L
- CC2420 communication modul (IEEE 802.15.4)
- Temperature sensor
- 4 LEDs
- Joystick
- Potentiometer



Motivation

Why an OS?

- Abstracts from hardware level
- Provide logical (device independent) interfaces
- Manage resources (CPU, memory, ...)



Sensor network operating system challenges

- Extreme resource constraints (processing, memory, energy, . . .)
- Different hardware
- Event driven – reactive to sensor inputs and communication
- Concurrently data processing
- Special operating systems needed



Examples of sensor network operating systems

- Tiny OS – <http://www.tinyos.net/>
- Contiki OS – <http://www.sics.se/contiki/>
- Mantis OS – <http://mantis.cs.colorado.edu>



TinyOS

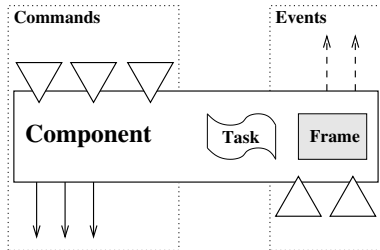
- Development since 2000 at the University of Berkley
- Open Source
- Objectives:
 - Minimal hardware requirements
 - Efficient parallel processing of data streams
 - Software modularity
- General structure
 - Scheduler
 - Schedules task in FIFO order
 - If task queue empty, puts the CPU into sleep mode
 - Components
- Programming language: NesC
 - C like syntax, ...
 - ... but more set up to the TinyOS model



TinyOS-Components I

Each component consists of four parts:

- Command handler
- Event handler
- Frame with static size
- one or more tasks



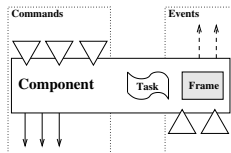
and can be realized in hardware as well as software.



TinyOS Components II

Commands

- Send by higher level components
- Return value shows failure or success



Events

- Triggered by hardware interrupts on the lowest level
- May trigger events on higher level components

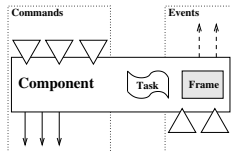
Commands and events are always yield small and fast actions without waiting times.



TinyOS Components III

Frames

- Data memory of the component
- Size is static and known at compile time



Tasks

- Doing the work
- Started by events, commands or other tasks of same component
- One at a time, only events may interrupt shortly
- Run to completion
 - + One stack is required for all tasks
 - Blocking or busy waiting not allowed



Contiki OS

- Adam Dunkel (Swedish Institute of Computer Science)
- Open Source
- Since March 2003 ported on 20 platforms
- System consists of a kernel and libraries
- Kernel:
 - Based on processes and events
 - Schedules and dispatches events to running processes
 - Provide CPU-multiplexing and message passing
 - Only platform independent code
 - Not multi-thread
- Libraries:
 - On top of kernel
 - Provide additional functionality often platform dependent code

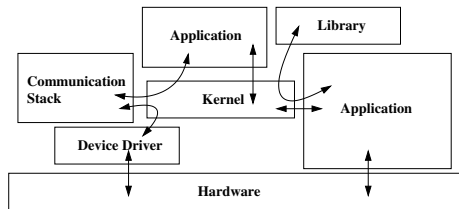


Contiki OS–Libraries

- Optionally linked to programs
- Run-time loading and linking possible
- Full access to underlying hardware
- Contain all platform dependent code

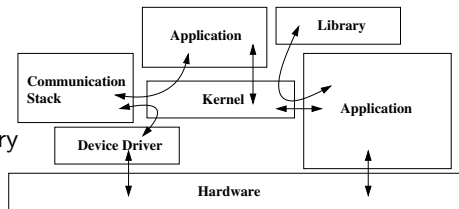
- Examples:

- Multi-threading
- Device driver
- Network stack (TCP/IP)
- Web browser
- GUI
- ...



Contiki OS–Processes

- Defined by an event handler function and ...
- ... an optional poll handler function
- Not preemptable and run to completion
- Fast context switches
- Extensions
 - Proto-threads
 - Multi-threading as a library



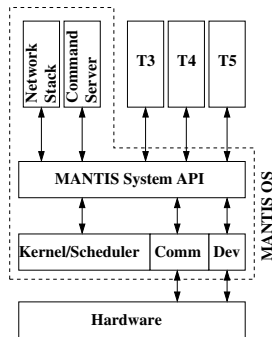
Contiki OS–Events

- Processes use events as a mean for IPC
- Events may be send to one or all processes
- Synchron:
 - Target process is immediately scheduled
 - After finishing, the caller gets control back
- Asynchron:
 - Event is queued
 - Target process is scheduled later
- Energy management has to be done by the application itself
 - Kernel provides only the fill level of the event queue



MANTIS OS – Multimodal NeTworks of In-situ Sensors

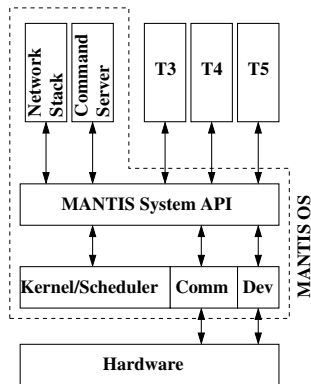
- Development since 2003 at the University of Colorado
- Open Source
- Objective: as simple as possible use and programming
- Features:
 - Classical layered architecture
 - Traditional stack based preemptive multi-threading model
 - I/O synchronisation via MUTEX
 - Network stack
 - Device driver
 - Kernel and API plain C



MANTIS OS

Threads

- Subset of POSIX-Threads
- Kernel manages RAM dynamically and allocate space on thread creation
- Threads are managed in a table with max. 12 entries
- Idle-thread is used for energy management if all other threads are blocked



Scheduler

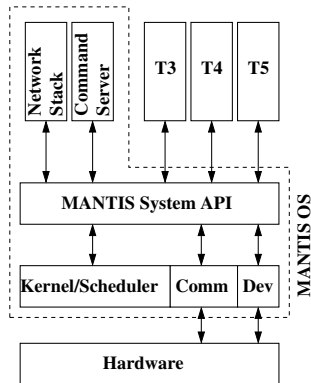
- UNIX-Style
- Priority based round-robin thread scheduling



MANTIS OS

Network Stack

- Four layer architecture:
 - Application
 - Network
 - MAC
 - Physical layer
- Layers are implemented as one or more threads
- Standard API between layers
- Flexible concept
- No TCP/IP



Comparison of the OS – I

| | TinyOS | Contiki | MANTIS OS |
|-----------------|-------------------------------|---|------------------------------|
| Target Platform | Sensor node | MiniOS not only for sensor nodes | Sensor node |
| Structure | Components | Micro-Kernel, Libraries | Classical layer structure |
| Programming | nesC | C | C |
| Scheduling | FIFO | Event based, Proto-threads, Multi-Threading via libraries | Priority based round robin |
| Communication | Via self developed components | TCP/IP | Via self developed protocols |
| Energy saving | Scheduler | Application | Idle-thread |



Comparison of the OS – II

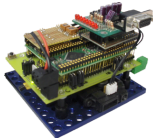
- Sizes of the operating systems are dependent on actual configuration.
- Values describe different applications in combination with the operating system and should give an impression what is achievable.

| | TinyOS | Contiki | MANTIS OS |
|-------------|------------|-----------|-----------|
| RAM Usage | 226 Bytes | 250 Bytes | 500 Bytes |
| Flash Usage | 3.5 kBytes | 4 kBytes | 14 kBytes |



Conclusion

- Three typical representatives of sensor network OS presented
- There are others like FreeRTOS, SOS, ...
- Deployment application context dependent



References I



Adam Dunkels and Björn Grönvall and Thiemo Voigt.

Contiki - a lightweight and flexible operating system for tiny networked sensors. In *Proceedings of the First IEEE Workshop on Embedded Networked Sensors (Emnets-I)*, Tampa, Florida, USA, November 2004.



Shah Bhatti, James Carlson, Hui Dai, Jing Deng, Jeff Rose, Anmol Sheth, Brian Shucker, Charles Gruenwald, Adam Torgerson, and Richard Han.

Mantis os: an embedded multithreaded operating system for wireless micro sensor platforms.

Mob. Netw. Appl., 10(4):563–579, 2005.



CHIPCON.

<http://www.chipcon.com>.



Crossbow.

<http://www.xbow.com>.



Adam Dunkels, Oliver Schmidt, Thiemo Voigt, and Muneeb Ali.

Protothreads: Simplifying event-driven programming of memory-constrained embedded systems.

In *Proceedings of the Fourth ACM Conference on Embedded Networked Sensor Systems (SenSys 2006)*, Boulder, Colorado, USA, November 2006.



References II



Jason Hill, Robert Szewczyk, Alec Woo, Seth Hollar, David Culler, and Kristofer Pister.

System architecture directions for networked sensors.
SIGPLAN Not., 35(11):93–104, 2000.



SOWNET.

<http://www.sownet.nl>.

