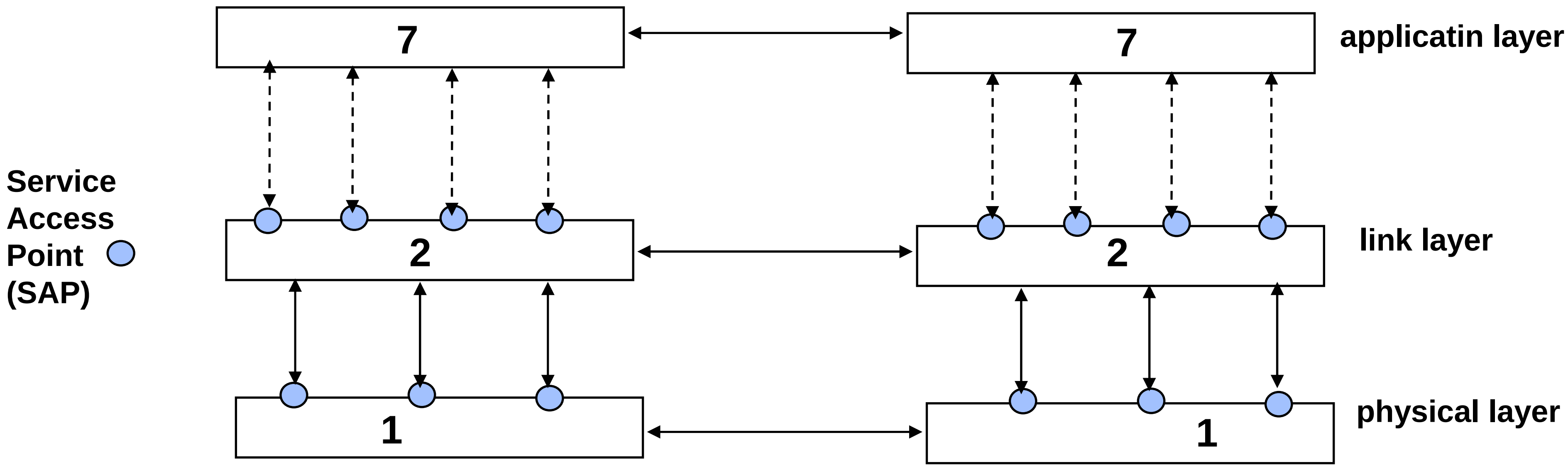
Physical Layer

Common Layering in the fieldbus area



- Assumption: homogeneous, closed system

- Not all layers are necessary (e.g. routing)
  Empty layers in the ISO/OSI- model
- Higher layers directly access the SAPs of lower layers.

- Efficiency improvement
- Direct mapping of layer 7 services to layer 2 functionality.

The Physical Layer Issues

- Asynchronous serial transmission   (character oriented)
- Synchronous serial transmission   (bitsynchronization)

- Bit coding:
  - NRZ (Non-Return-to-Zero
  - Manchester Code
  - MFM (Modified-Frequency-Modulation)

- Modulation and data transmission:
  - Base band — Example: Morsetel. / Ethernet
  - Broad band — Example: Radio, TV, Cabel-TV, Modem
    Modulation: AM, FM

- Transmissionmedia:
  - Fiber — (Multi-Mode, Single-Mode)
  - Copper — (Twisted Pair, Coaxial)
  - Radio — (Frequency band)
  - Satellite — (Geostationary, orbiting)

Properties of communication networks

Constraining factors:

- **Transfer rate, (capacity, bandwidth)**
- **Propagation latency**

Transfer rates:

Morse-telegraph: < 100Bit/sec
Telegraphy: < 150 Bit/sec
Phone: ~ 50Kbit /sec
Serial RS232: ~ 100Kbit/sec
Field bus: few Kbit/sec ... ~ 1Mbit/sec
Ethernet: 10-1000Mbit/sec
High speed networks: >> Gbit/sec

Latency: **Satellite connection (2 x 35700 km): ~ 240 ms,**
cabel (trans-atlantic) (~ 6.000 km): ~ 20 ms

Topology: point-to-point, star, bus, tree, grid, multi-level....

How much information can be transfered over a line?

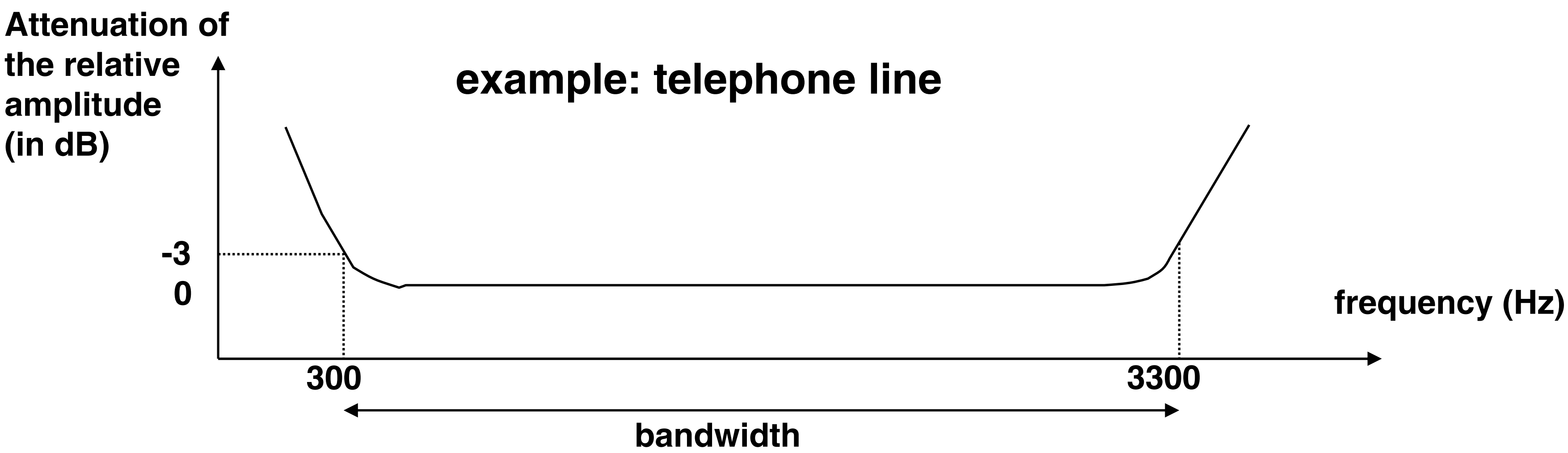
limiting factors:

→ **bandwidth of the channel**

→ **noise**

→ **The bandwidth limits the number of transitions, i.e. the frequency of switching from one signal level to the other**

→ **Noise limits the ability to distinguish between multiple signal levels**



Capacity of a channel (Shannon):

$$C = B \bullet ld\,(P_s + P_n) / P_n = B \bullet ld\,(1 + P_s / P_n)$$

C : capacity of a channel   (measured in Bit/sec (bps))
P_s : signal strength (measured in µW, mW, W)
P_n : noise (measured in µW, mW, W)
B: bandwidth
P_s / P_n : signal-to-noise ratio (dB) = $20 \bullet \log_{10} (P_s/P_n)$

Example:
Telephone: bandwidth 3000Hz, signal-to-noise ratio 60 dB
(corresponds to a relation 1000/1)

$C = 3000 \bullet ld\,(1+1000) = 3000 \bullet 9{,}97 = 29900$ Bit/sec (bps)

U_1/U_0	dB	comment
1000	60	amplification
100	40	amplification
10	20	amplification
3,16	10	amplification
2	6	amplification
1,414	3	amplification
1	0	(1:1) transmission
0,7071	-3	attenuation
0,5	-6	attenuation
0,316	-10	attenuation
0,1	-20	attenuation
0,01	-40	attenuation
0,001	-60	attenuation

$$L = 20 \cdot \lg \frac{U_1}{U_0}$$

Bps and BAUD

Bps (Bit/sec) defines a Bit rate

BAUD defines the number of level transitions

Bit/sec is contraint by the channel capacity !

BAUD is constraint by the bandwidth !

Basic methods to increase the bps-Rate at a given BAUD rate of the channel:

- **distinguish multiple levels**
- **Coding with the smallest number level transitions**

Coding options (base band)

- level
- pulse width
- transistions

Bit coding:

NRZ	(Non-Return-To-Zero)
Manchester	
MFM	(Modified-Frequency-Modulation)

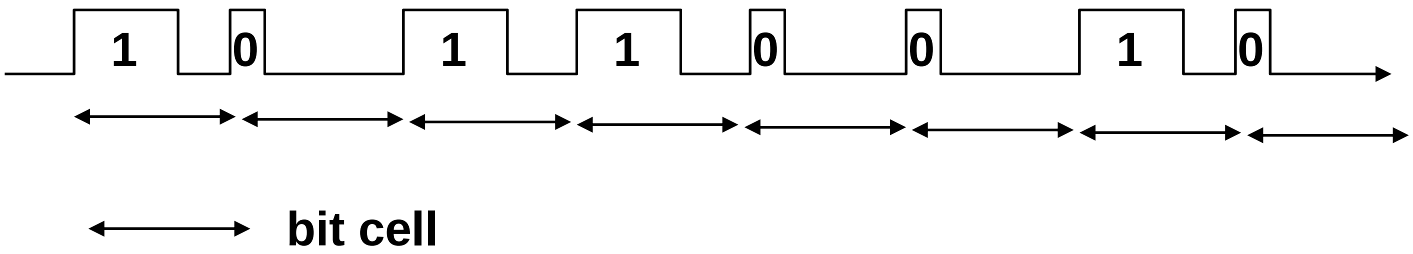
Problems:

- synchronization
- number of transitions
- constant/variable frame length

A (bad) example

RZ: (Always) Return to Zero (PWM)

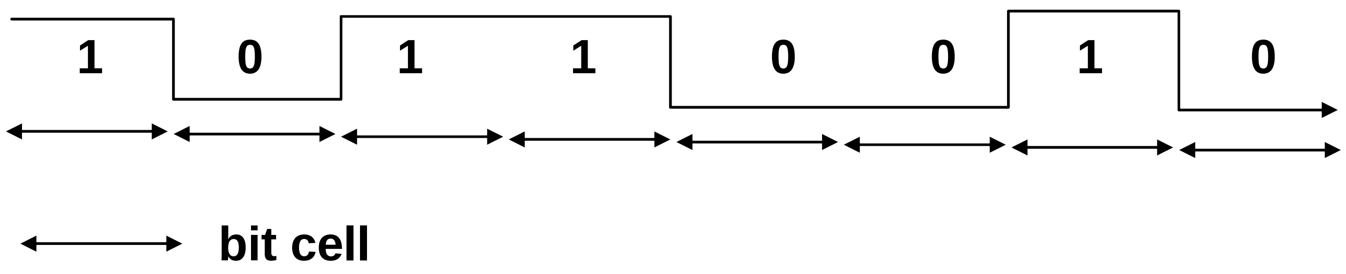
Example: 1011 0010



NRZ Codes

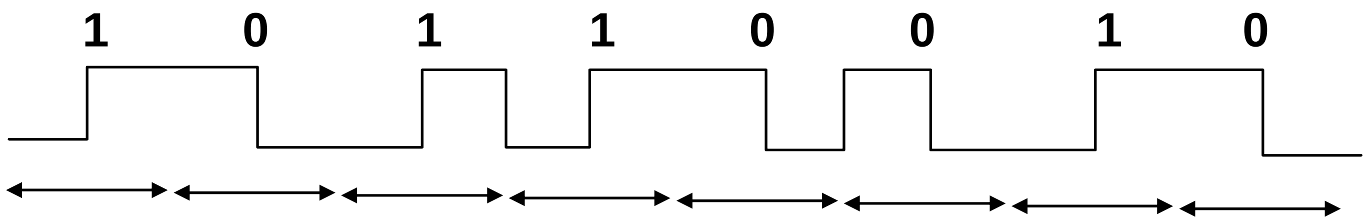
NRZ: Non Return to Zero

Example: 1011 0010



Manchester Coding

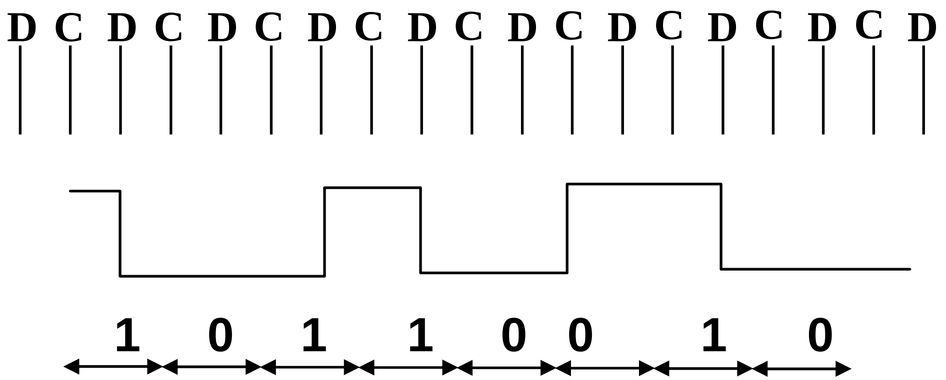
Example: 1011 0010



bit cell

MFM (Modified Frequency Modulation)

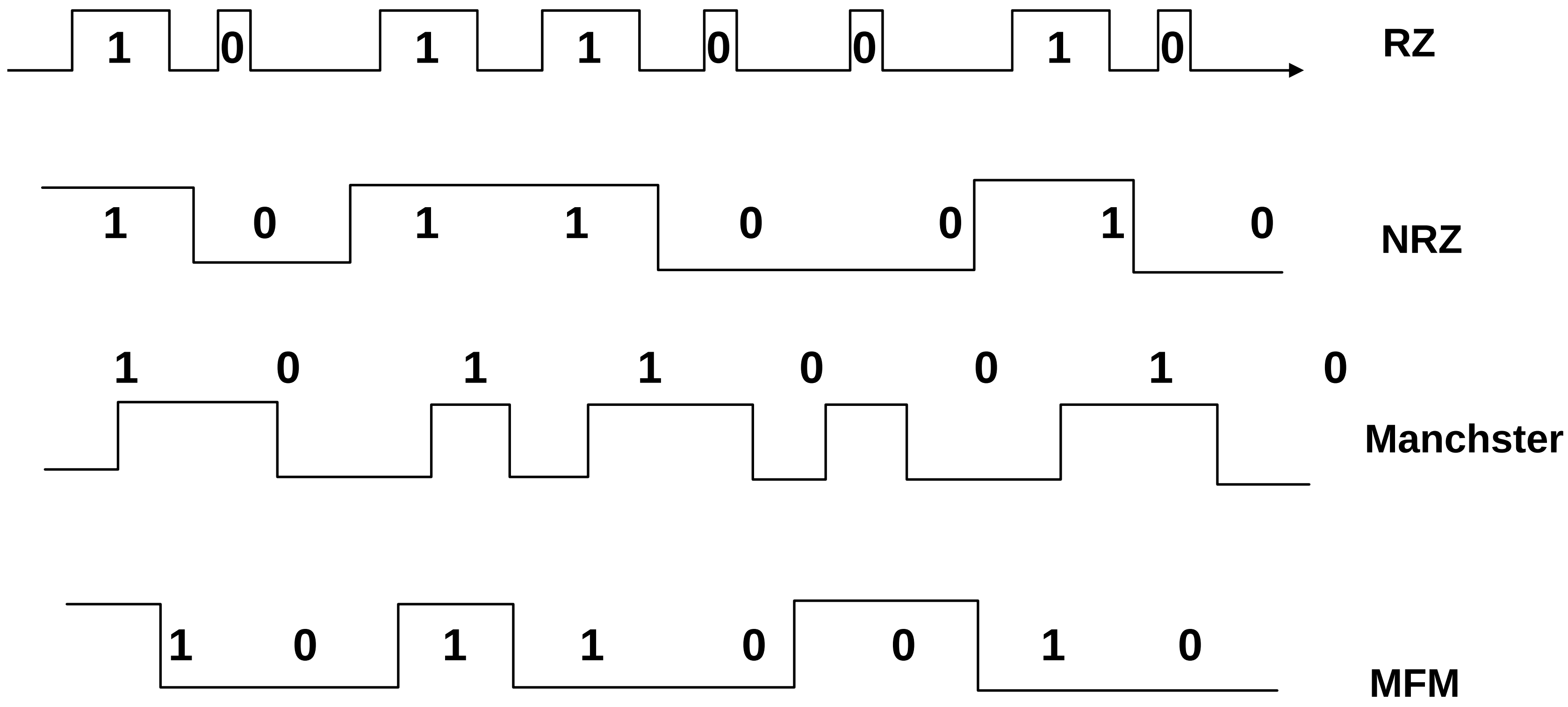
Example: 1011 0010



1: always trasition at a data point (D)
0: no transition at a data point (D)

multiple consecutive "0" : transistion at a clock point (C)

Comparison of Codes



Comparison of Codes

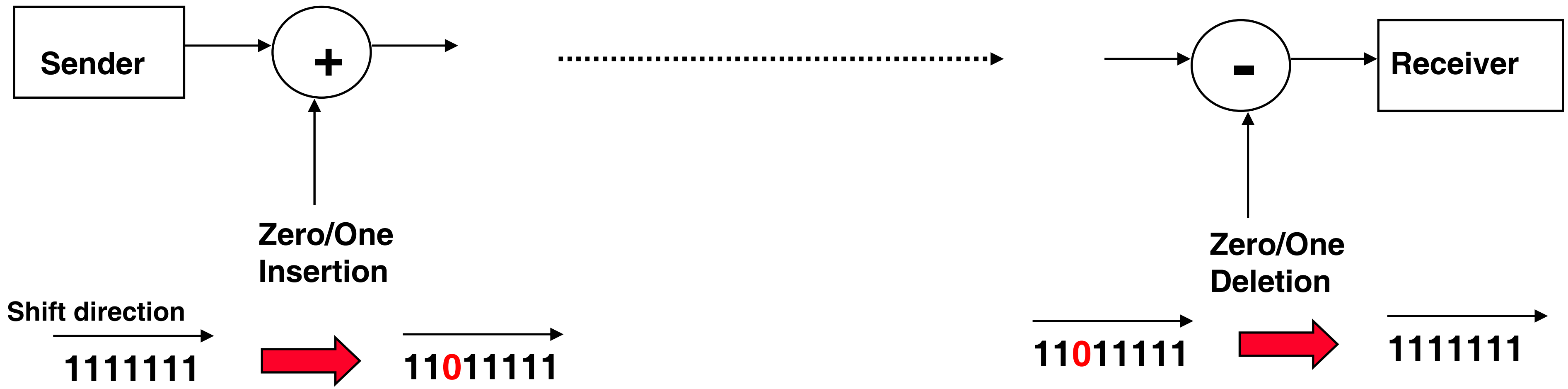
Type	Synchronization	transitions/Bit average/max		fixed length
RZ	Y	2	2	Y
NRZ	N	>0,5	1	Y
NRZ*	Y	>0,5	1	N
Manchester	Y	1,5	2	Y
MFM	Y	>0,5	1	Y

NRZ*: NRZ mit Bit Stuffing

Bit-Stuffing

When a long sequence of identical values "0" or "1" occurs, bit stuffing inserts a complemetary signal level after a fixed specified number of equal signal levels.

Sender transparently inserts stuff bits. The receiver re-establishes the original message by removing the respective stuff bits.



Bit stuffing to identify message boundaries

Example: HDLC (High Level Data Link Control)

Problem:
In character-oriented protocols control and seperation characters (STX, ETX, etc.) can be identified.
In bit-oriented protocols any combination of bits as data is possible.

➡ **How to identify control information?**

HDLC-Frame

Flag	address	control	info		flag
0111 1110	8	8	optional	FCS	0111 1110

control →

I-Frame - Information Frame: data transport
S-Frame - Supervisor Frame: flow control, e.g. ACK, re-transmission
U-Frame - Un-numbered Frame: additional control info e.g. connect, disconnect

How to distinguish the control info 01111110 from data ?

Bit stuffing to identify message boundaries

goal: recognizing the flag "01111110".

method: The sender normally inserts a stuff bit "0" after 5 consecutive "1". Therefore there is a max. number of five consecutive "1". The flag is inserted AFTER the bist stuffing stage in the sender and detected and removed before the receiver stuffing stage.